1.1.athe address of the image of process can changed in main memory

b. the start address

c. because we need to use the fregmant

1.2 small page size: : can run more programs with the same memory

more page fault, bigger page table, the content of page is small

1,3 paging is fixed size by hardware.

1.4segment has variable page size. segmant is used to store important process

2.every part that changed the data  
b.sem.value—and sem.value ++ is infact a inconsistency

c. no it is a dead lock.

3. 0.98\*0.2+0.2\*20=4.196

4.

|  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|  | 3 | 3 | 3 | 1 | 1 | 1 | 1 | 1 | 1 |
|  |  | 2 | 2 | 2 | 2 | 2 | 3 | 6 | 7 |

2.

|  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
|  | 3 | 3 | 3 | 1 | 1 | 1 | 1 | 1 | 1 |
|  |  | 2 | 2 | 2 | 2 | 2 | 3 | 6 | 7 |

.

5.no need to go to main memory and secondary memory

b. maybe cache is not big enough

6.a/ Does not hinder the performance of processes and hence results in greater system throughput.

b/The page fault ratio of a process can not be solely controlled by the process itself. The pages in memory for a process depends on the paging behavior of other processes as well.

7.a/good b/bad c/bad

8. if there are too many files. it will be slow

9. small page size: : can run more programs with the same memory

more page fault, bigger page table, the content of page is small

10. segment has variable page size. segmant is used to store important process

11. cause they are lighter than process. The thread shares the same virtual memory space. Process is in the memory.